

Towards Trace Metrics via Functor Lifting

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My Talk Today

Introduction

Lifting Functors [FSTTCS '14]

Behavioral (Pseudo-)Metrics [FSTTCS '14]

Trace Metrics via Functor Lifting

Final Remarks

Introduction

Behavioral Equivalence vs Behavioral Metric

Behavioral Equivalence

Two systems behave **exactly** the same.

(bisimilarity, trace/language equivalence, failure equivalence ...)

Behavioral (Pseudo-)Metrics

Two systems behave **similarly**, i.e. they have a small **distance**.

(bisimilarity **pseudometric**, trace/language **pseudometric** ...)

Our Central Contributions

- ▶ FSTTCS '14: Behavioral pseudometrics for coalgebras of Set-functors.
- ▶ This talk: Trace pseudometrics (using Eilenberg-Moore).

Examples of Trace Distances

- ▶ NFA $L_x, L_y : A^* \rightarrow 2$

$$d_{2^{A^*}}(L_x, L_y) = c^{\min\{n \in \mathbb{N}_0 \mid \exists w \in A^n. L_x(w) \neq L_y(w)\}}$$

where $c \in (0, 1)$.

- ▶ Probabilistic Moore automata $p_x, p_y : A^* \rightarrow [0, 1]$

$$d_{[0,1]^{A^*}}(p_x, p_y) = c_1 \cdot \sum_{w \in A^*} \left(\frac{c_2}{|A|} \right)^{|w|} |p_x(w) - p_y(w)| .$$

where $c_1, c_2 \in (0, 1]$, $c_1 + c_2 \leq 1$.

Lifting Functors [FSTTCS '14]

Definition: Lifting of a Functor (to \mathbf{PMet})

\bar{F} is a *lifting* of F if the diagram commutes.

$$\begin{array}{ccc} \bar{F} & \xrightarrow{\quad} & \mathbf{PMet} \\ & & \downarrow \mathcal{U} \\ F & \xrightarrow{\quad} & \mathbf{Set} \end{array}$$

In Plain Terms: Our Task

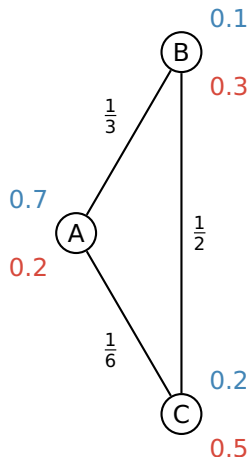
Given: $F: \mathbf{Set} \rightarrow \mathbf{Set}$ and a pseudometric space (X, d)

Define: pseudometric d^F on FX

Pseudometrics on Probability Distributions

Example Based on Optimal Transportation [Villani, 2009]

$$X = \{A, B, C\}, \quad d_X: X \times X \rightarrow [0, 1], \quad \mu, \nu: X \rightarrow [0, 1]$$



0.1 Transportation Problem

- ▶ A, B, C: producers/consumers
- ▶ supply: μ , demand: ν (in %)
- ▶ transporting along d costs d
- ▶ minimize total cost of transport

Optimal Solution

- ▶ first use local production (cost 0)
- ▶ transport 0.2 from A to B: $\frac{1}{3} \cdot 0.2 = \frac{1}{15}$
- ▶ transport 0.3 from A to C: $\frac{1}{6} \cdot 0.3 = \frac{1}{20}$

$$d(\mu, \nu) = \frac{1}{15} + \frac{1}{20} = \frac{7}{60}$$

$$\mu, \nu \in \mathcal{D}(X), \quad t \in \mathcal{D}(X \times X), \quad t: \underbrace{X}_{\text{from}} \times \underbrace{X}_{\text{to}} \rightarrow \underbrace{[0, 1]}_{\text{percentage}}$$

supply is used up

$$\forall x \in X. \sum_{y \in X} t(x, y) = \mu(x)$$

$$\iff \mathcal{D}\pi_1(t) = \mu$$

demand is satisfied

$$\forall y \in X. \sum_{x \in X} t(x, y) = \nu(y)$$

$$\iff \mathcal{D}\pi_2(t) = \nu$$

Definition: Coupling

$F: \text{Set} \rightarrow \text{Set}$, $t_1, t_2 \in FX$. The set of **couplings** of t_1 and t_2 is

$$\Gamma_F(t_1, t_2) := \left\{ t \in F(X \times X) \mid \begin{array}{l} F\pi_1(t) = t_1 \\ F\pi_2(t) = t_2 \end{array} \right\}.$$

$$\text{costs}_d(t) = \sum_{(x,y) \in X \times X} t(x,y) \cdot d(x,y)$$

$$d: X \times X \rightarrow [0, 1] \quad \rightsquigarrow \quad \text{costs}_d: \mathcal{D}(X \times X) \rightarrow [0, 1]$$

Definition: Evaluation Functor & Evaluation Function

$$\begin{array}{ccc} \mathbf{Set}/[0, T] & \xrightarrow{\tilde{F}} & \mathbf{Set}/[0, T] \\ X \xrightarrow{f} [0, T] & \mapsto & FX \xrightarrow{Ff} \underbrace{F[0, T] \xrightarrow{\text{ev}_F} [0, T]}_{\text{evaluation function}} \end{array}$$

Definition: Wasserstein Distance

For $t_1, t_2 \in FX$ define

$$d^{\downarrow F}(t_1, t_2) = \inf \left\{ \tilde{F}d(t) \mid t \in \Gamma_F(t_1, t_2) \right\} .$$

Proposition: Wasserstein Pseudometric

$d^{\downarrow F}$ is a pseudometric if

- ▶ F preserves weak pullbacks
- ▶ ev_F satisfies some constraints (well-behaved)

Behavioral (Pseudo-)Metrics [FSTTCS '14]

Final Coalgebra and Behavior

The Essential Basics in a Nutshell (i.e. on one Slide)

Definition: Final Coalgebra

$\omega: \Omega \rightarrow F\Omega$ is final, if for all $c: X \rightarrow FX$ there is a unique coalgebra homomorphism $[\cdot]_c: X \rightarrow \Omega$.

$$\begin{array}{ccc} X & \xrightarrow{\forall c} & FX \\ \downarrow \exists! [\cdot]_c & & \downarrow F[\cdot]_c \\ \Omega & \xrightarrow{\omega} & F\Omega \end{array}$$

Definition: Behavioral Equivalence

$$\forall x, y \in X: \quad x \sim_c y \iff [x]_c = [y]_c$$

Definition: Behavioral Pseudometric [FSTTCS '14]

$$\forall x, y \in X: \quad d_c(x, y) = d_\omega([x]_c, [y]_c)$$

for d_ω on Ω turning ω into a final \bar{F} -coalgebra

For **all** coalgebras $c: X \rightarrow FX$ we want to have:

$$d_c(x, y) = d_\omega([\![x]\!]_c, [\![y]\!]_c). \quad (1)$$

Especially for $c = \omega$, where ω is an iso and $[\![\cdot]\!]_\omega = \text{id}_\Omega$.

Definition: Behavioral Pseudometric

Define $d_c: X \times X \rightarrow [0, \top]$ as least fixpoint of:

$$d_c(x, y) = d_c^F(c(x), c(y))$$

Theorem: Behavioral Pseudometric

1. $c = \omega$ yields final \bar{F} -coalgebra.
2. If \bar{F} preserves isometries, (1) holds for all c .
3. If ω is obtained via the final chain construction, \bar{F} preserves isometries and metrics, then d_ω is a proper metric.

Trace Metrics via Functor Lifting

Coalgebraic Traces in Eilenberg-Moore

Our coalgebras are of the shape

$$c: X \rightarrow FTX$$

where T is a monad on \mathbf{Set} and $F = B \times _{}^A$.

- ▶ A nondeterministic automaton is a \mathbf{Set} -coalgebra

$$c: X \rightarrow 2 \times (\mathcal{P}_{\text{fin}} X)^A.$$

- ▶ A probabilistic Moore automaton is a \mathbf{Set} -coalgebra

$$c: X \rightarrow [0, 1] \times (\mathcal{D}X)^A.$$

Known: To obtain trace semantics via final coalgebra, we need to move to the Eilenberg-Moore category of T .

Definition: Lifting of a Functor to $\mathcal{EM}(T)$

\hat{F} is a lifting of F if the diagram commutes.

$$\begin{array}{ccc} \hat{F} & \hookrightarrow & \mathcal{EM}(T) \\ & & \downarrow u^T \\ F & \hookrightarrow & C \end{array}$$

Proposition: Liftings and Distributive Laws

$$\hat{F}: \mathcal{EM}(T) \rightarrow \mathcal{EM}(T) \quad \xleftrightarrow{\cong} \quad \lambda: TF \Rightarrow FT$$

Given λ and an FT -coalgebra $c: X \rightarrow FTX$ define:

$$c^\# := \left(TX \xrightarrow{Tc} TFTX \xrightarrow{\lambda_{TX}} FTTX \xrightarrow{F\mu_X} FTX \right)$$

This F -coalgebra yields a \hat{F} -coalgebra $\mu_X \rightarrow \hat{F}\mu_X$.

Traces in Eilenberg-Moore

via the generalized powerset construction [Silva et al., 2013, Jacobs et al., 2015]

$$\begin{array}{ccccc} X & \xrightarrow{\eta_X} & TX & \xrightarrow{\llbracket \cdot \rrbracket_{c^\sharp}} & \Omega \\ c \downarrow & & \swarrow c^\sharp & & \downarrow \omega \\ FTX & \xrightarrow{F\llbracket \cdot \rrbracket_{c^\sharp}} & & & F\Omega \end{array}$$

Definition: Trace Equivalence

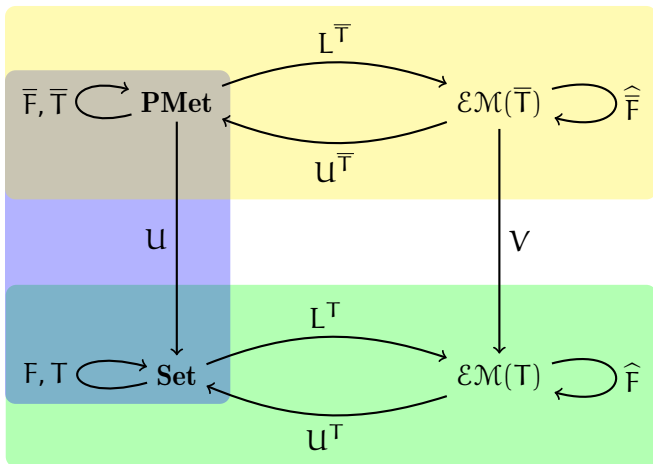
$$\forall x, y \in X: \quad x \approx_c y \iff \llbracket \eta(x) \rrbracket_{c^\sharp} = \llbracket \eta(y) \rrbracket_{c^\sharp}$$

Definition: Trace Pseudometric [new! 😊]

$$\forall x, y \in X: \quad d_c(x, y) = d_\omega\left(\llbracket \eta(x) \rrbracket_{c^\sharp}, \llbracket \eta(y) \rrbracket_{c^\sharp}\right)$$

for d_ω on Ω turning ω into a final \bar{F} -coalgebra

The “Big Picture”



Our Results

... Ensure the Validity of the Aforementioned Approach

- ▶ compositionality

$$\overline{F} \overline{G} = \overline{FG}$$

- ▶ lifting of natural transformations

$$\lambda: F \Rightarrow G \quad \rightsquigarrow \quad \overline{\lambda}: \overline{F} \Rightarrow \overline{G}$$

- ▶ lifting of monads

$$(T, \eta, \mu) \quad \rightsquigarrow \quad (\overline{T}, \overline{\eta}, \overline{\mu})$$

- ▶ we obtain the **trace pseudometrics** of our initial examples!

Final Remarks

Summary, Future and Related Work

Our Contribution

Trace **pseudometrics** via functor lifting and determinization!

Future/Related Work

- ▶ What about coalgebras $X \rightarrow TFX$ in $\mathcal{Kl}(T)$?
- ▶ Our lifting approach is clearly inspired by [van Breugel and Worrell, 2005]. What about a **logical characterization** of distances as in [Desharnais et al., 1999]?
- ▶ **Directed** metrics (simulation distances) [de Alfaro et al., 2009]
- ▶ Vary metric on $[0, \top]$ [Chatzikokolakis et al., 2014]
- ▶ **Algorithms** (polynomial-time [Chen et al., 2012], on-the-fly [Bacci et al., 2013], . . .)
- ▶ A **fibrational** perspective (work in progress)

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